

Functional Design

Principles, Patterns, and Practices

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Functional Design

```
for (Driver d : drivers) {
    if (!d.getRumors().equals(rumors))
      return true;
    }
    return false;
}
```

A quick perusal of this code will convince you that it is written in a very traditional OO style and that the objects encapsulate their own state relatively well.

CLOJURE

When writing the Clojure version I did not start out with a design sketch. Rather, I depended upon my TDD tests to help me with the design. The tests are as follows:

```
(ns gossiping-bus-drivers-clojure.core-spec
 (:require [speclj.core :refer :all]
            [gossiping-bus-drivers-clojure.core :refer :all]))
(describe "gossiping bus drivers"
 (it "drives one bus at one stop"
    (let [driver (make-driver "d1" [:s1] #{:r1}4)
         world [driver]
         new-world (drive world)]
      (should= 1 (count new-world))
      (should= :s1 (-> new-world first :route first))))
 (it "drives one bus at two stops"
    (let [driver (make-driver "d1" [:s1 :s2] #{:r1})
         world [driver]
         new-world (drive world)]
      (should= 1 (count new-world))
      (should= :s2 (-> new-world first :route first))))
```

^{4. #{. . .}} represents a set in Clojure. A set is a list of items that has no duplicates.

```
(it "drives two buses at some stops"
  (let [d1 (make-driver "d1" [:s1 :s2] #{:r1})
        d2 (make-driver "d2" [:s1 :s3 :s2] #{:r2})
        world [d1 d2]
        new-1 (drive world)
        new-2 (drive new-1)]
    (should= 2 (count new-1))
    (should= :s2 (-> new-1 first :route first))
    (should= :s3 (-> new-1 second :route first))
    (should= 2 (count new-2))
    (should= :s1 (-> new-2 first :route first))
    (should= :s2 (-> new-2 second :route first))))
(it "gets stops"
  (let [drivers #{{:name "d1" :route [:s1]}
                  {:name "d2" :route [:s1]}
                  {:name "d3" :route [:s2]}}]
    (should= {:s1 [{:name "d1" :route [:s1]}
                   {:name "d2" :route [:s1]}]
              :s2 [{:name "d3", :route [:s2]}]}
             (get-stops drivers)))
  )
(it "merges rumors"
  (should= [{:name "d1" :rumors #{:r2 :r1}}
            {:name "d2" :rumors #{:r2 :r1}}]
           (merge-rumors [{:name "d1" :rumors #{:r1}})
                          {:name "d2" :rumors #{:r2}}])))
(it "shares gossip when drivers are at same stop"
  (let [d1 (make-driver "d1" [:s1 :s2] #{:r1})
        d2 (make-driver "d2" [:s1 :s2] #{:r2})
        world [d1 d2]
        new-world (drive world)]
    (should= 2 (count new-world))
    (should= #{:r1 :r2} (-> new-world first :rumors))
    (should= #{:r1 :r2} (-> new-world second :rumors))))
```

There are some interesting similarities between the Java tests and the Clojure tests. They are both quite wordy; although the Clojure tests contain half as many lines. The Java version has 12 tests whereas the Clojure version has only 8. This difference has a lot to do with the way the two different solutions were partitioned. The Clojure tests also play pretty fast and loose with the data.

Consider, for example, the "merges rumors" test. The merge-rumors function expects a list of drivers; however, the test does not create completely formed drivers. Rather, it creates abbreviated structures that look like drivers as far as the merge-rumors function is concerned.

The solution is all contained in a single, very short file:

```
(ns gossiping-bus-drivers-clojure.core
   (:require [clojure.set :as set]))
(defn make-driver [name route rumors]
   (assoc<sup>5</sup> {} :name name :route (cycle<sup>6</sup> route) :rumors rumors))
(defn move-driver [driver]
```

^{5.} assoc adds elements to a map. (assoc {} :a 1) returns {:a 1}.

^{6.} cycle returns a lazy (and "infinite") list that simply repeats the input list endlessly. Thus, (cycle [1 2 3]) returns [1 2 3 1 2 3 1 2 3 ...].

```
(update driver :route rest))
(defn move-drivers [world]
  (map move-driver world))
(defn get-stops [world]
 (loop [world world
         stops {}]
    (if (empty? world)
      stops
      (let [driver (first world)
            stop (first (:route driver))
            stops (update stops stop conj driver)]
        (recur (rest world) stops)))))
(defn merge-rumors [drivers]
 (let [rumors (map :rumors drivers)
        all-rumors (apply set/union<sup>8</sup> rumors)]
      (map #(assoc % :rumors all-rumors) drivers)))
(defn spread-rumors [world]
 (let [stops-with-drivers (get-stops world)
        drivers-by-stop (vals9 stops-with-drivers)]
    (flatten¹0 (map merge-rumors drivers-by-stop))))
(defn drive [world]
 (-> world move-drivers spread-rumors))
(defn drive-till-all-rumors-spread [world]
 (loop [world (drive world)
```

^{7.} The update function returns a new map with one element changed. (update m k f a) changes the k element of m by applying the function (f e a), where e is the old value of element k. Thus, (update {:x 1} :x inc) returns {:x 2}.

^{8.} The union function is from the set namespace. Notice the ns at the top aliases the clojure.set namespace to just set.

^{9.} vals returns a list of all the values in a map. keys returns a list of all the keys in a map.

^{10.} The flatten function turns a list of lists into a list of all the elements. So (flatten [[1 2][3 4]]) returns [1 2 3 4].

```
time 1]
(cond
  (> time 480) :never
  (apply = (map :rumors world)) time
  :else (recur (drive world) (inc time)))))
```

This solution is 42 lines, whereas the Java solution is 145 lines spread among five files.

Both solutions have the concept of a Driver, but I made no attempt to encapsulate the concepts of Route, Stop, and Rumor into independent objects. They all just happily live within the Driver.

Worse, the Driver "object" is not an object in the traditional OO sense. It has no methods. There is one method in the system, move-driver, that operates on a single Driver, but it's just a little helper function for the more interesting move-drivers function.

Six out of the eight functions take only the world as an argument. Thus, we might say that the only true object in this system is the world, and it has five methods. But even that is a stretch.

Even if we decide that the Driver is a kind of object, it is not mutable. The simulated world is nothing more than a list of immutable Drivers. The drive function accepts the world and produces a new world in which all the Drivers have been moved one step, and Rumors have been spread at any stop where more than one Driver has arrived.

That drive function is an example of an important concept. Notice how the world passes through a pipeline of functions. In this case there are only two, move-drivers and spread-rumors, but in larger systems the pipeline can be quite long. At each stage along that pipeline the world is modified into a slightly new form.

This tells us that the partitioning of this system is not about objects, but about functions. The relatively unpartitioned data passes from one independent function to the next.

You might argue that the Java code is relatively straightforward, whereas the Clojure code is too dense and obscure. Believe me when I say that it does not take very long to get comfortable with that density and that the perceived obscuration is an illusion based on unfamiliarity.

Is the lack of partitioning in the Clojure version a problem? Not at its current size; but if this program were to grow the way most systems grow, that problem would assert itself with a vengeance. Partitioning OO programs is a bit more natural than partitioning functional programs because the dividing lines are much more obvious and pronounced.

On the other hand, the dividing lines we chose for the Java version are not guaranteed to lead to an *effective* partitioning. The warning is in the drive function of the Clojure program. It seems likely that a better partitioning of this system might lie along the different operations that manipulate the world, rather than things like Routes, Stops, and Rumors.

CONCLUSION

We saw some differences in the Prime Factors and Bowling Game katas; but the differences were relatively minor. The differences in the Gossiping Bus Drivers kata were much more pronounced. This is likely because that last kata was a bit larger than the first two (I'd say twice the size), and also because it was a true finite state machine.

A *finite state machine* moves from state to state, taking actions that depend upon the incoming events and the current state. When such systems are written in an OO style, the state tends to be stored in mutable objects that have dedicated methods. But in a functional style, the state remains